

Seminar
Database and Information Systems

Security Primitives and Methods

Betreuer: Boris Stumm

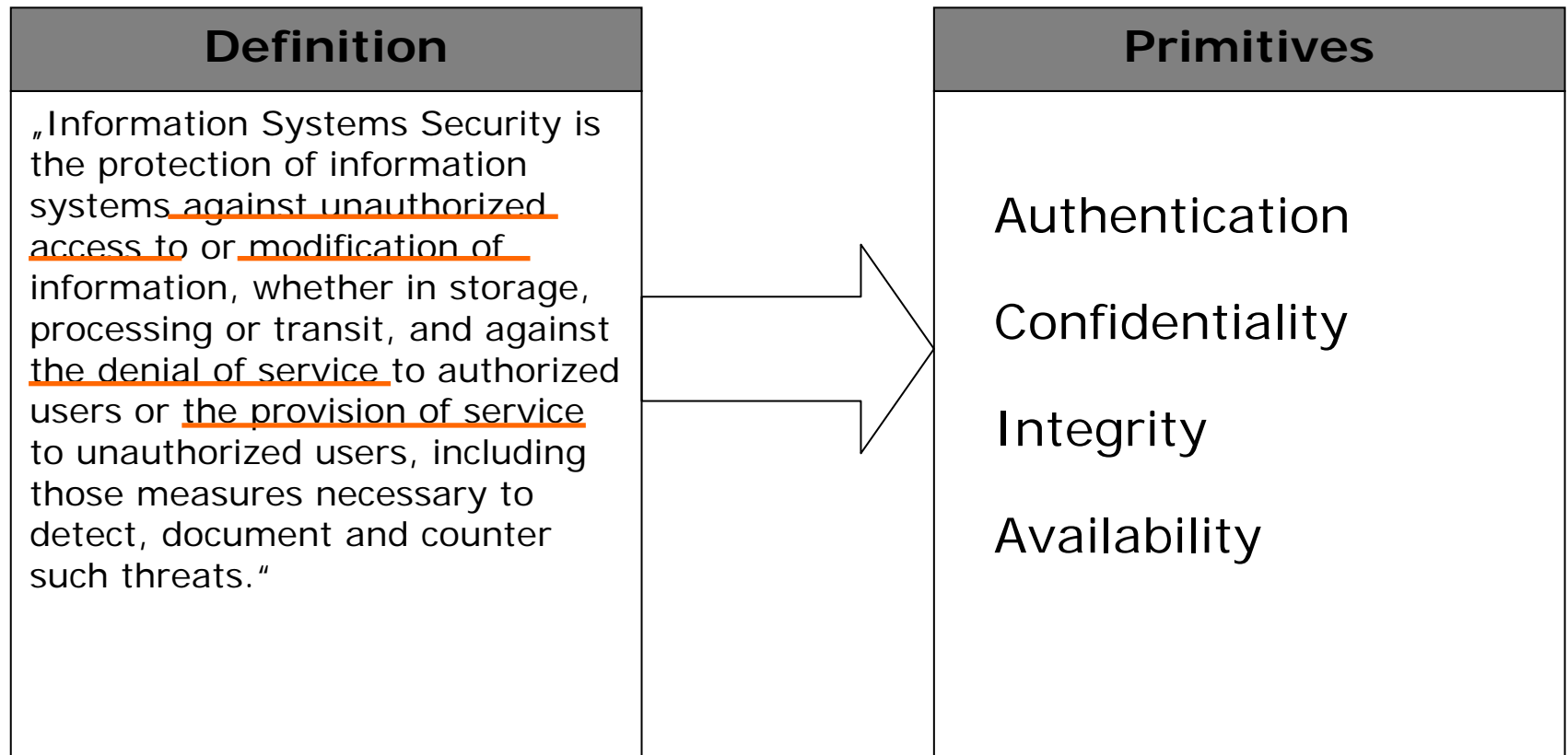
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Introduction

- Security Primitives
- Methods and Applications
- Protocols
- Conclusion

Security primitives derived from the definition of information systems security



Realization of security primitives

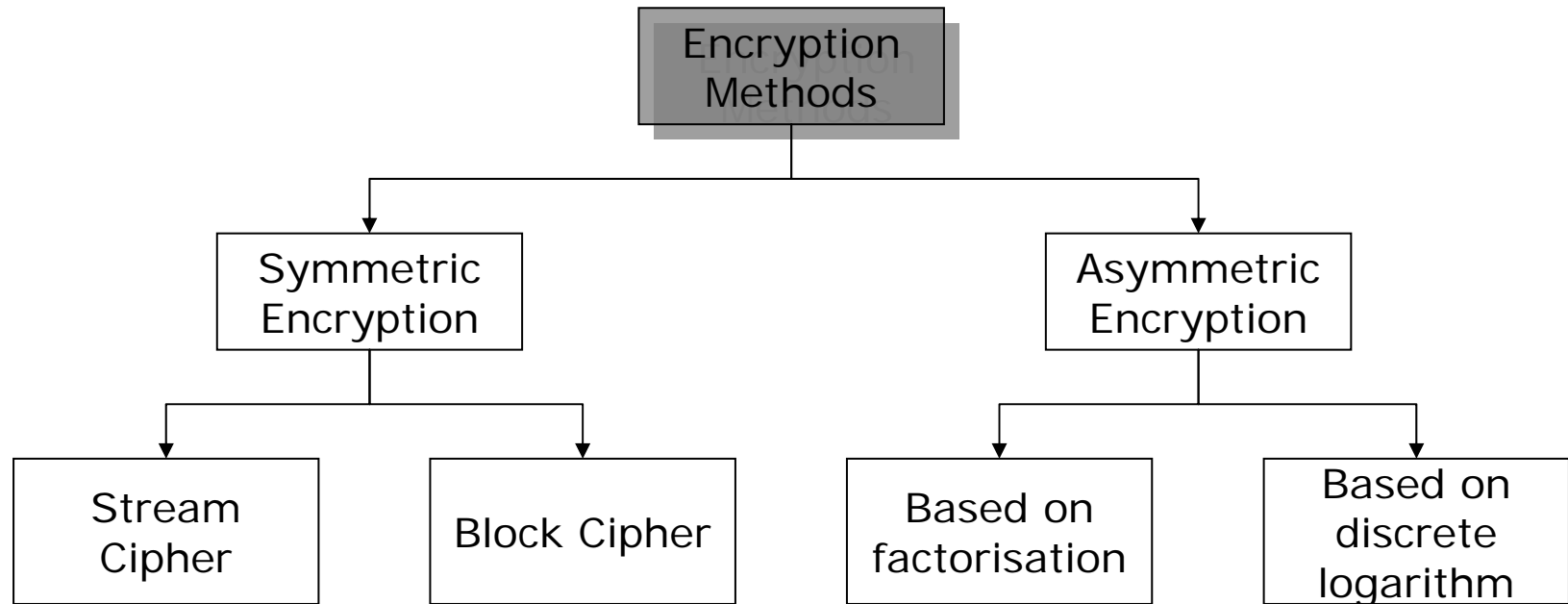
- Confidentiality
 - Encryptions
- Integrity
 - Hash Function, Digital Signature
- Authentication
 - Digital Signature
- Availability
 - backups, fail-over system, disaster recovery plans

Cryptography is the basic technique

Cryptography = kryptos (hidden or secret) + graphos (writing)



Categories of encryption methods



Symmetric Encryption

- Uses the same key to encipher and decipher.
- Two fundamental classes:
 - Stream Cipher
 - Encrypts the plaintext letter by letter
 - Block Cipher
 - First divides the plaintext into blocks, then enciphers

Stream Cipher: Vigenère Cipher

- The formula of Vigenère Cipher is mathematically represented as:
Ciphertext = Plaintext + Key modulo 26

- Here is an example:

Plaintext: tob eor not tob eth ati sth equ est ion
 Key: run run run run run run run run run run run
 Ciphertext: kio vie eig kio vnu rnv jnu vkh vmg zia

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
a	b	c	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z

Stream Cipher: One-Time Pad

- The only secure encryption method.
 - The key length is as long as the plaintext.
 - The key is not repeatedly used.
- It is expensive for most of the applications.
- How to exchange the key safely?

Block Cipher: Playfair

Plaintext:

The rain in Spain stays mainly on the plain.

Step 1:

theraininspainstaysmainlyontheplain

Step 2:

th er ai ni ns pa in st ay sm ai nl yo nt he pl ai n

Step 3:

th er ai ni ns pa in st ay sm ai nl yo nt he pl ai nz

Step 4:

rk pn sw su rt al us to mw oa sw te mf ro up am sw ge

P	A	L	M	E
R	S	T	O	N
B	C	D	F	G
H	I	K	Q	U
V	W	X	Y	Z

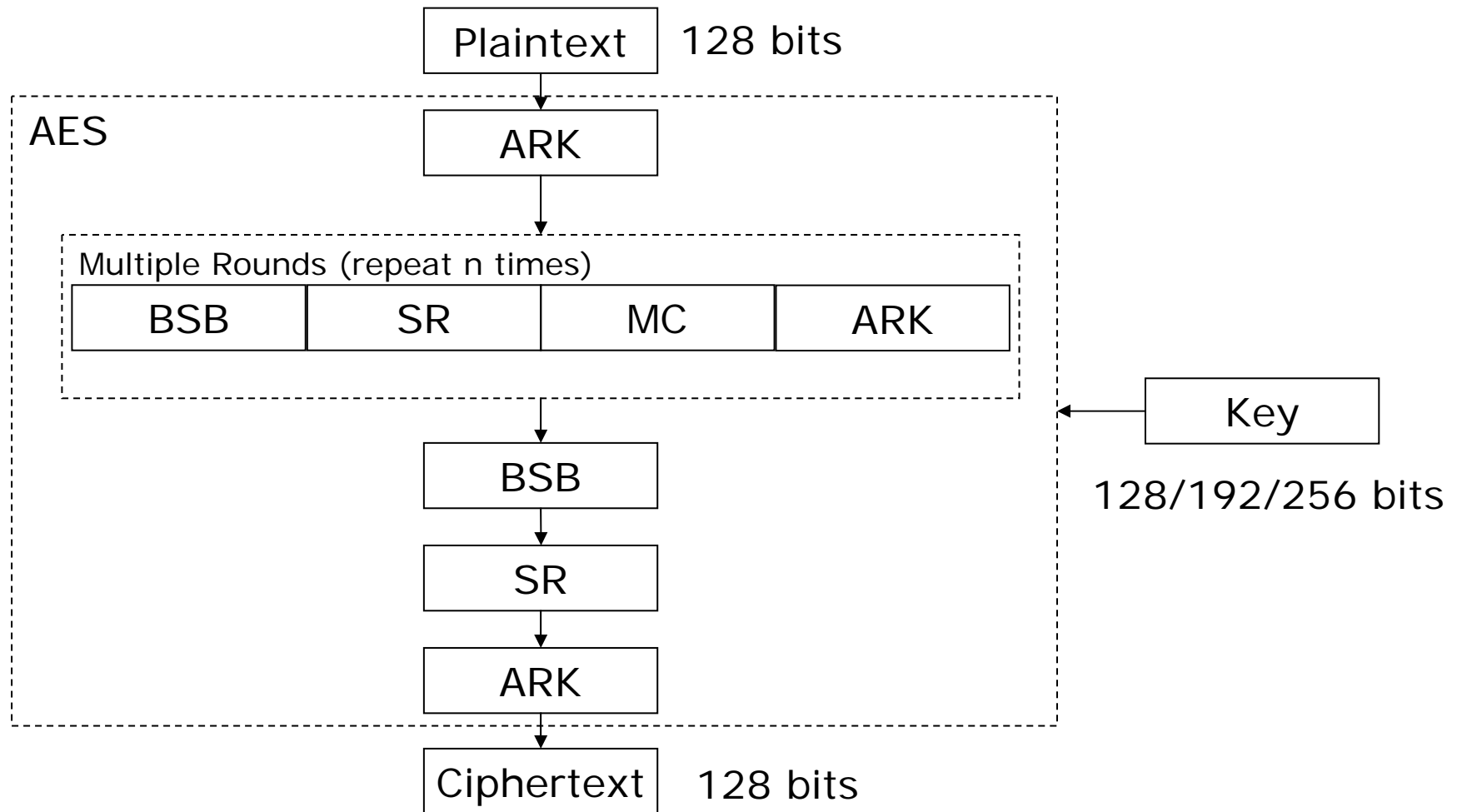
Block Cipher: AES

- In 1997 in America, the National Institute of Standards and Technology (NIST) initiated a process to select a symmetric-key encryption algorithm to be used to protect sensitive Federal information.
- In 2000 the algorithm invented by Vincent Rijmen and Joan Daemen was selected as Advanced Encryption Standard (AES)

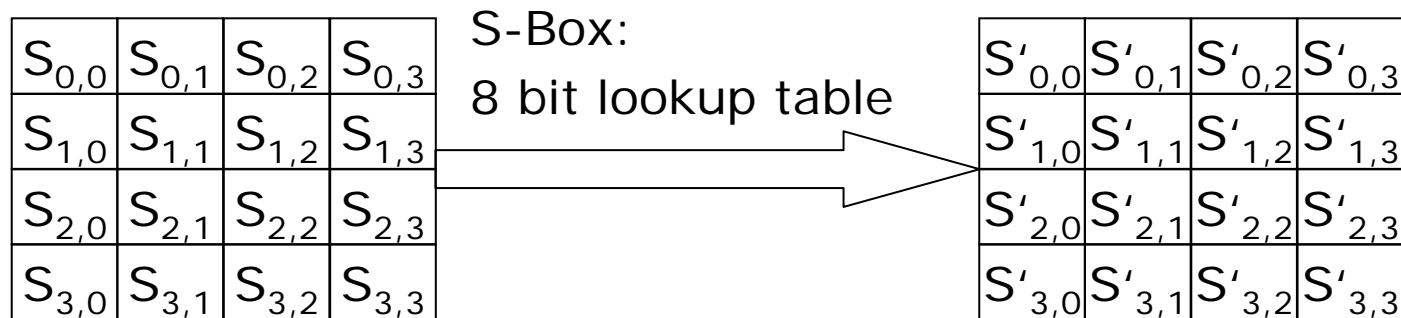
Basic data of AES

- AES has a block size of 128 bits, key size of 128, 192, 256 bits.
- The general working sectors of AES are:
 - Add Round Key (ARK);
 - Byte Substitution Byte (BSB);
 - Shift Row (SR);
 - Mix Column (MC)
- The number of rounds depends on the key size.
 - 128 bits key: 9 rounds
 - 192 bits key: 11 rounds
 - 256 bits key: 13 rounds

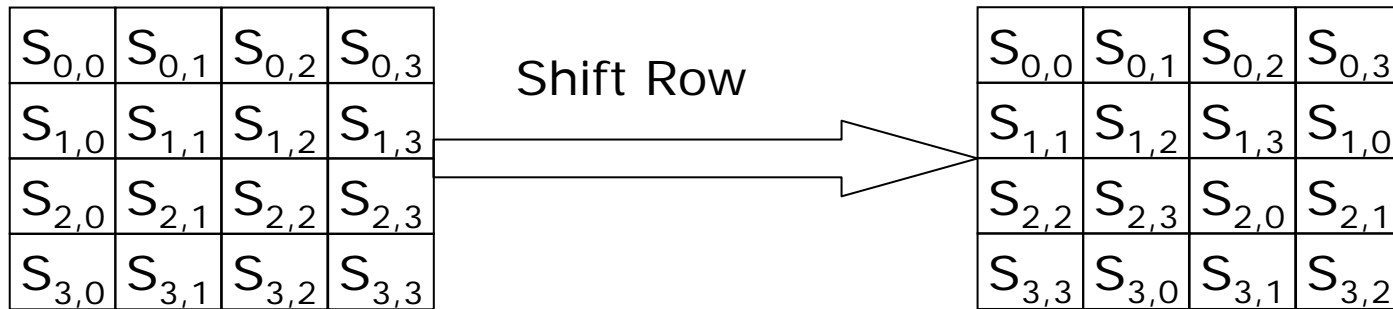
Working process of AES



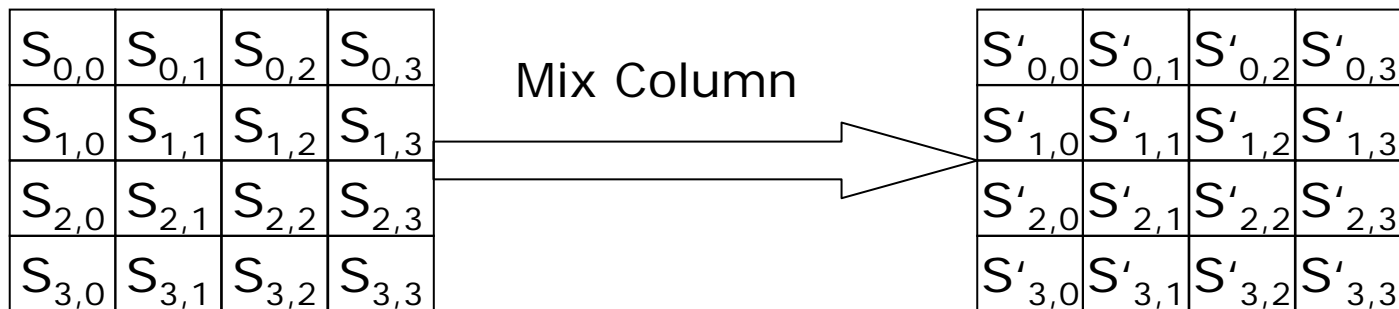
Byte Substitution Byte: Each byte will be replaced by another byte using the S-box



Shift Row: Shift the row according to the row number

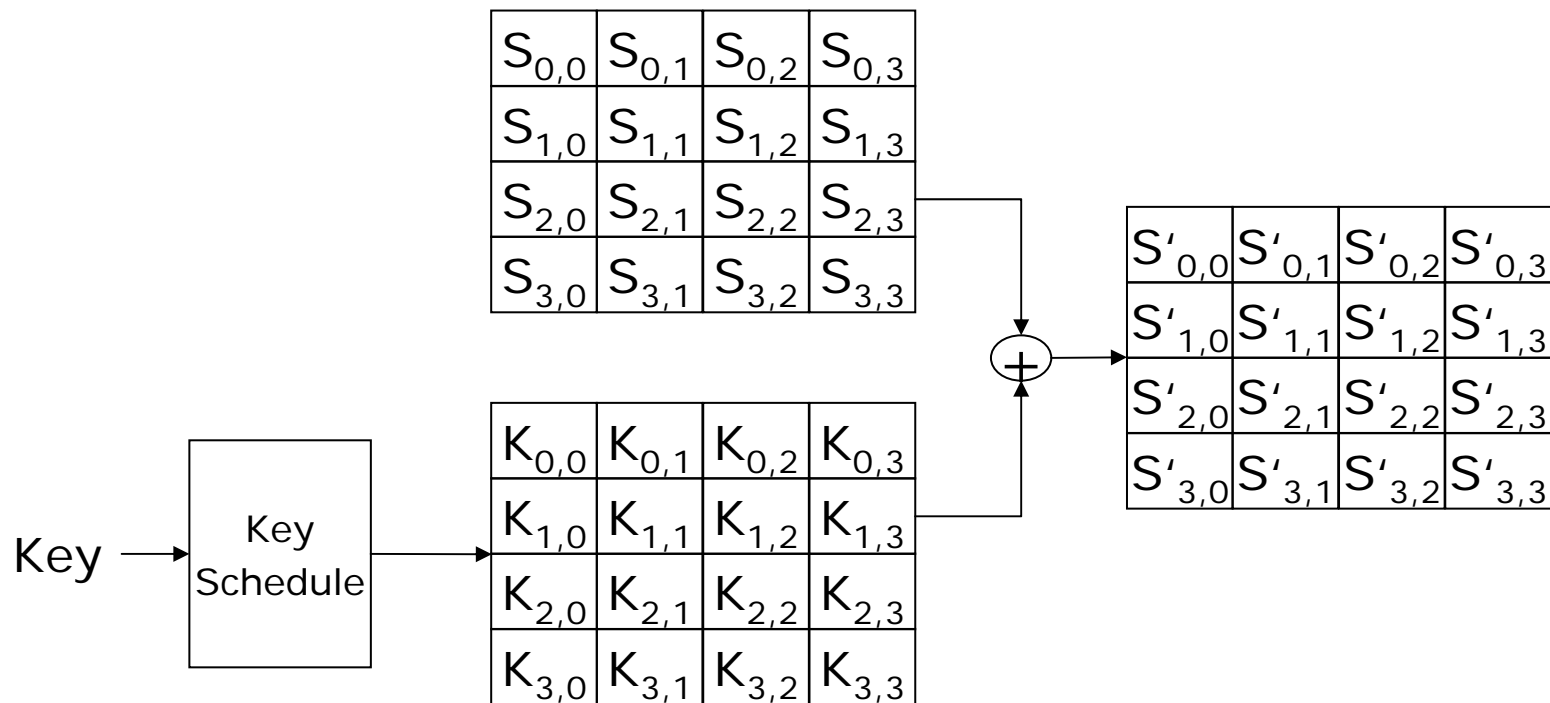


Mix Column: Every byte in a column will be mixed



$$\begin{pmatrix} S'_{3,0} \\ S'_{2,0} \\ S'_{1,0} \\ S'_{0,0} \end{pmatrix} = \begin{pmatrix} 02 & 01 & 01 & 03 \\ 03 & 02 & 01 & 01 \\ 01 & 03 & 02 & 01 \\ 01 & 01 & 03 & 02 \end{pmatrix} \begin{pmatrix} S_{3,0} \\ S_{2,0} \\ S_{1,0} \\ S_{0,0} \end{pmatrix}$$

Add Round Key: The selected round key is added to the text with XOR



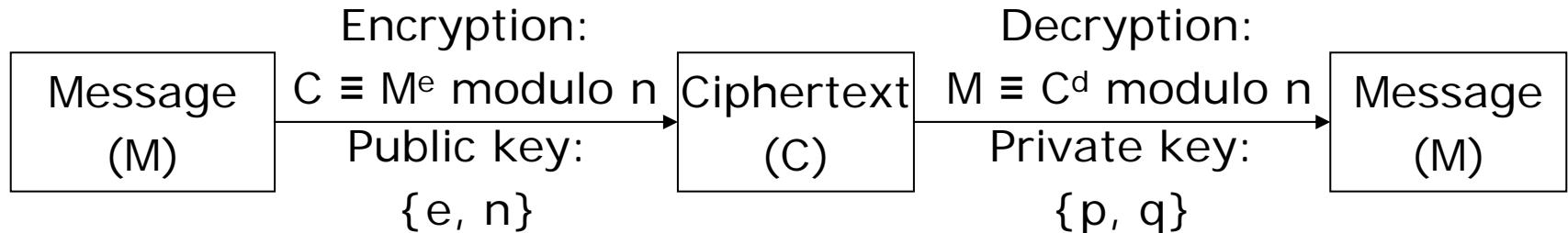
Asymmetric Encryption

- Uses different keys to encrypt and decrypt
 - Encryption: public key
 - Decryption: private key
- Based on the difficulty of solving a certain mathematical problem:
 - Factorisation: It is hard to solve a composite number into its factors if this number is big enough.
 - RSA
 - Discrete logarithm: For each natural number y , it can be expressed as $y \equiv g^x \text{ modulo } p$.
 - ElGamal

RSA: Background Knowledge

- RSA was invented by Ron Rivest, Adi Shamir and Len Adleman in 1978.
- The two basic mathematic theories supporting RSA are:
 - Fermat's Little Theorem:
For all primes p not dividing a , $a^{p-1} \equiv 1$ modulo p
 - Euler's function $f(n)$:
 $f(n)$ is the number of positive integers less than n with which it has no divisor in common.
So if n is the product of two primes p and q , then:
 $f(n) = (p-1)(q-1)$

RSA: encryption and decryption formulas



$$C^d \equiv \{M^e\}^d \equiv M^{ed} \equiv M^{1+k \cdot f(n)} \equiv M \cdot M^{k \cdot f(n)} \equiv M \times 1 \equiv M \text{ modulo } n$$

n : a natural number with factors p and q

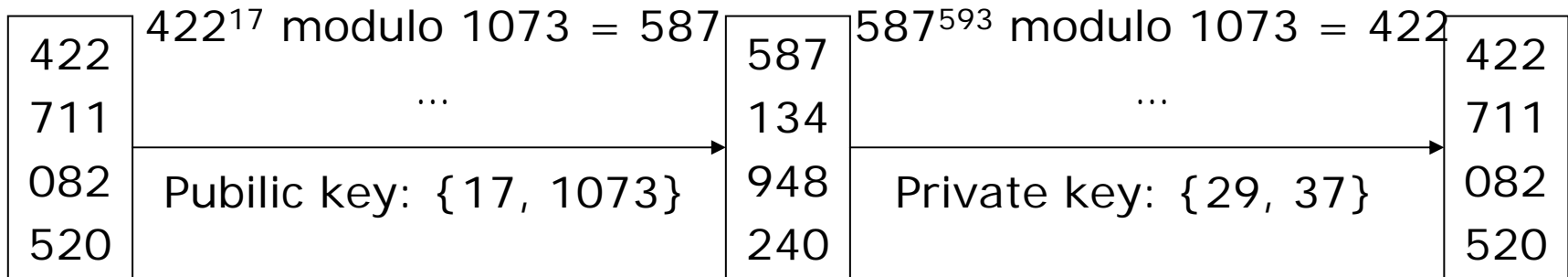
p : a prime number between $[0, n-1]$

q : a prime number between $[0, n-1]$

e : exponent between $[0, n-1]$ without common factors of $p-1$ and $q-1$

d : a number between $[0, n-1]$ with $d \cdot e \equiv 1 \text{ modulo } f(n)$

RSA: Example



$M = 422711082520;$ $n = 1073;$ $p = 29;$
 $q = 37;$ $e = 17;$ $d = 573;$
 $C = 587134948240$

ElGamal: Background

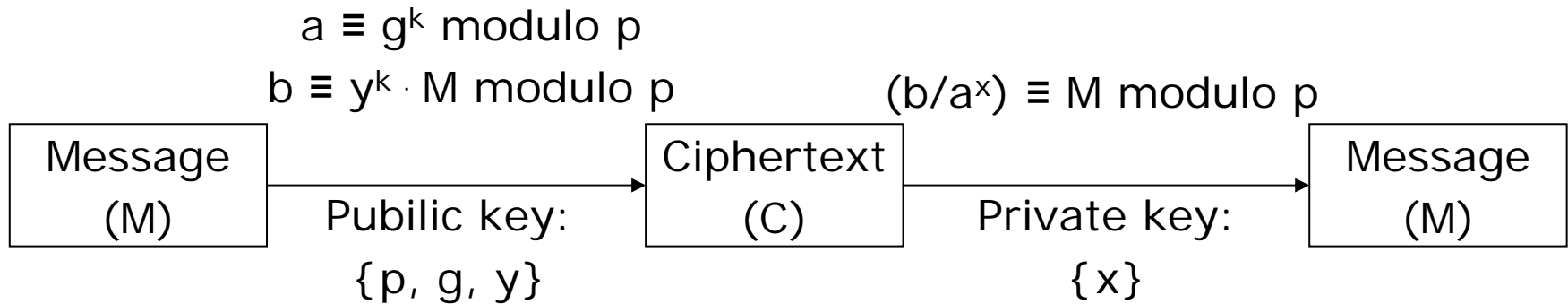
- ElGamal was invented by Taher Elgamal.
- The theory behind ElGamal is:
 - For each natural number y , there is a number x , which satisfies the equation: $y \equiv g^x \pmod{p}$. p is a random prime number. y, g, x are natural numbers smaller than p .
 - It is easy to find the corresponding x for y by calculation as long as p is small. If the number is big it will be very difficult.

Backup 2

- For example: suppose $p = 11$, $g = 7$:

y	x	because
1	10	$7^{10} \text{ modulo } 11 = 1$
2	3	$7^3 \text{ modulo } 11 = 2$
3	4	$7^4 \text{ modulo } 11 = 3$
4	6	$7^6 \text{ modulo } 11 = 4$
5	2	$7^2 \text{ modulo } 11 = 5$
6	7	$7^7 \text{ modulo } 11 = 6$
7	1	$7^1 \text{ modulo } 11 = 7$
8	9	$7^9 \text{ modulo } 11 = 8$
9	8	$7^8 \text{ modulo } 11 = 9$
10	5	$7^5 \text{ modulo } 11 = 10$

ElGamal: encryption and decryption formulas



x : a natural number between $[0, p-1]$

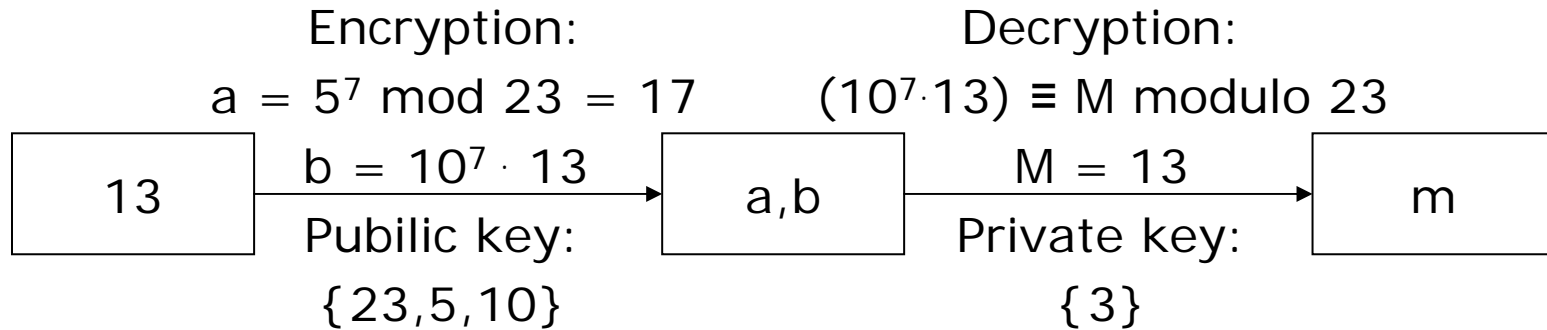
y : a natural number between $[0, p-1]$, $y \equiv g^x \text{ modulo } p$

g : a natural number between $[0, p-1]$

p : a random prime number

k : a random natural number

ElGamal: example



$$x = 3;$$

$$y = 10;$$

$$g = 5;$$

$$k = 7;$$

$$p = 23;$$

Comparison of Symmetric and Asymmetric Encryption

Symmetric

(+) The calculation in symmetric encryption is relatively easy, like addition, multiplication

(-) A large number of keys are needed

2 Persons: 1 keys

100 Persons: 4950 keys

Can be compensated by the application of Key Distribution Center (KDC).

(-) Risk to exchange the secret key safely.

Asymmetric

(+) No need to exchange the private key.

(+) The number of keys needed is reduced.

2 Persons: 2 keys

100 Persons: 200 keys

(-) The calculation is relatively complex.

(-) Comparing with symmetric encryption, the key size is big.

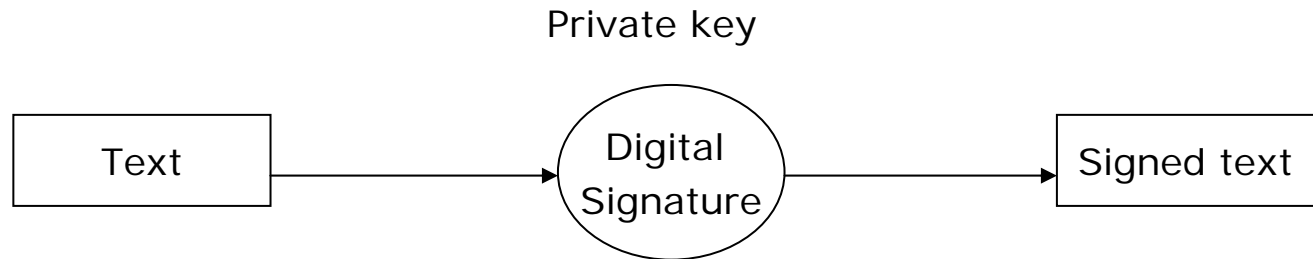
(-) It is possible to find the some mathematic solutions in the algorithm.

One-Way Hash Function

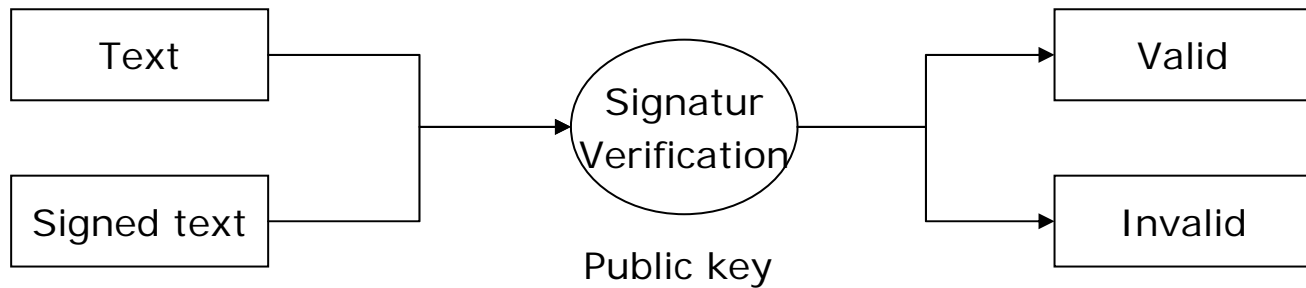
- One-way hash function should meet the conditions:
 - The hash value should be relatively small for any input
 - Two different messages should not have the same hash value
 - Only one-direction calculation is possible
- Can be used to create a MIC (message integrity code) to check integrity
- Can be used to compute message digests.

Digital Signature:

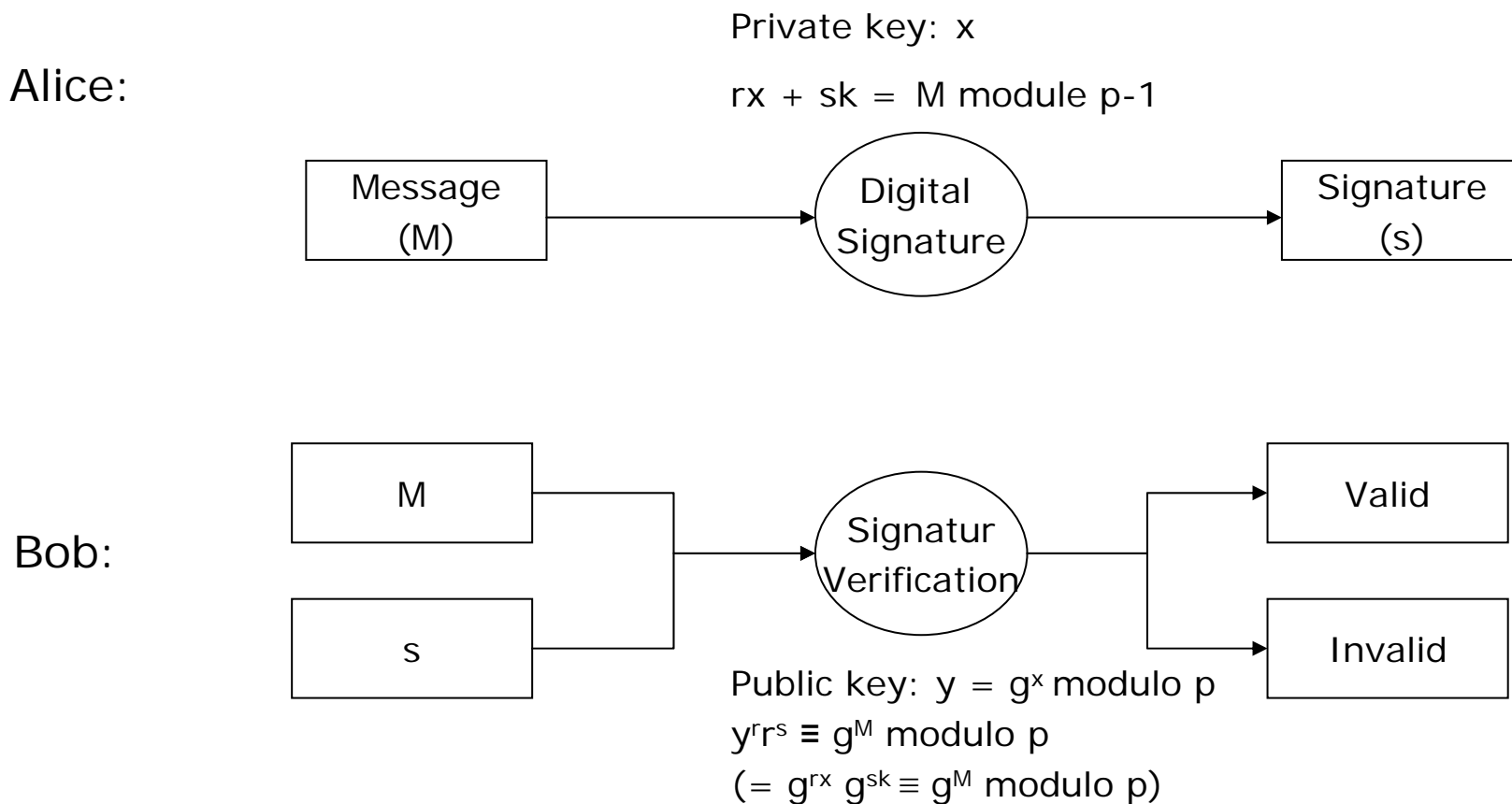
Sender:



Receiver:



Digital Signature: Using ElGamal Algorithm



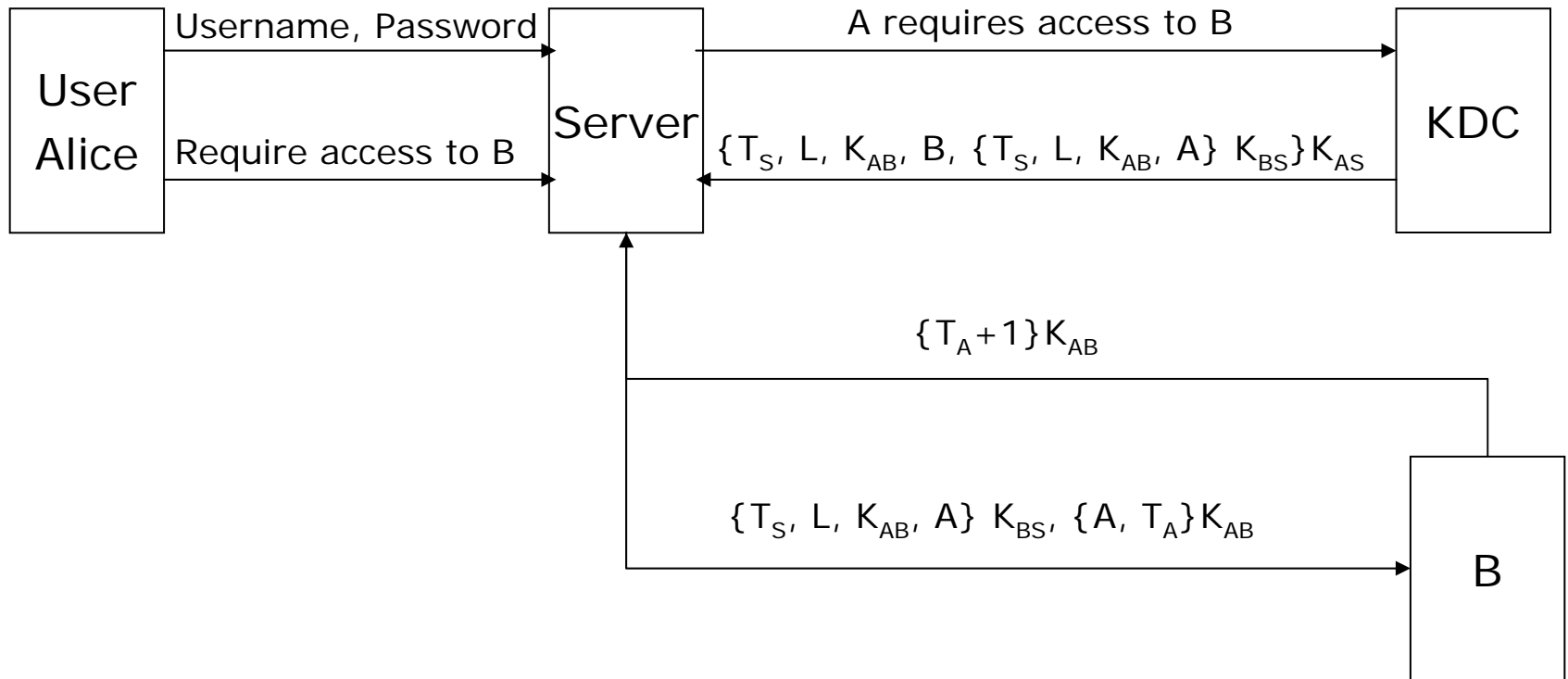
k, g are random natural numbers; p is a prime number; $r = g^k \text{ modulo } p$

Kerberos

- Kerberos is an authentication protocol using symmetric encryption.
- The name Kerberos comes from Greek mythology: Kerberos is the three-headed dog that guarded the entrance to Hades.
- In computer security, three heads represent client, server and KDC (key distribution center).

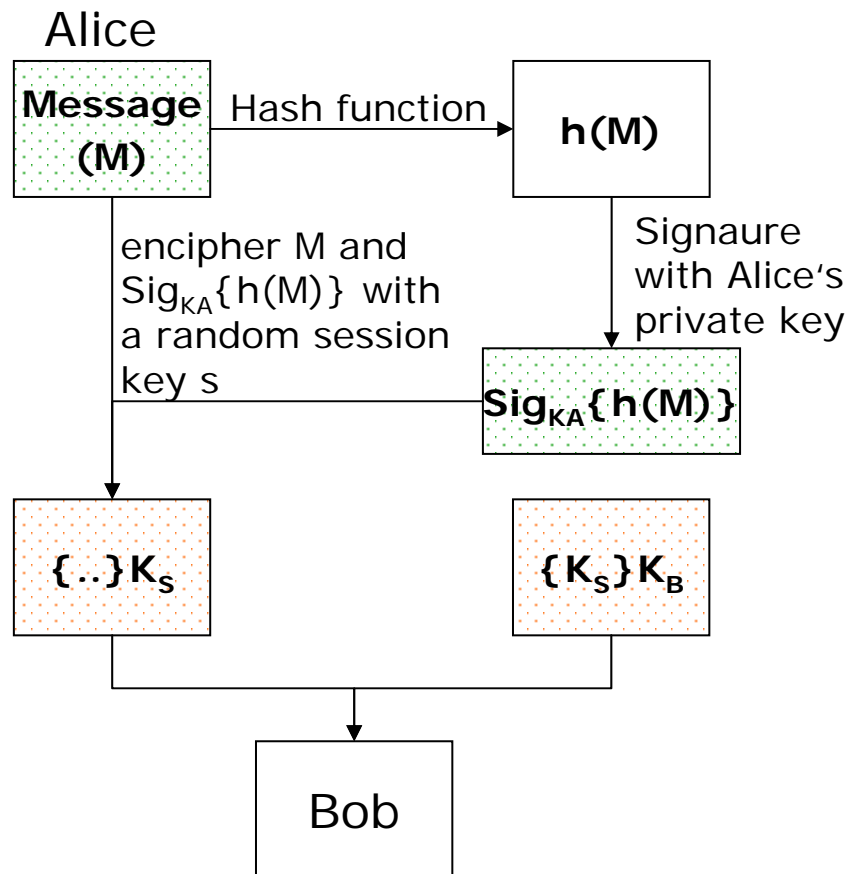


Kerberos



PGP (Pretty Good Privacy)

- PGP is a secure mail protocol. It was invented by Phil Zimmermann in 1991.
- PGP uses both public key encryption and private key encryption.
- PGP is used in Evolution, Eudora, Mozilla Thunderbird.
- Plug-ins implementing PGP is also available for Outlook Express.



Conclusion

- Security Primitives
 - Confidentiality, Authentication, Integrity, Availability
- Methods and Protocols
 - Encryption, Hash Function, Digital Signature, Kerberos, PGP
- Security is only a relative concept