

Chapter 9 – Web Services Coordination and Transactions



Coordination - Motivation

- Interactions are typically more complex than simple invocations
- Need to coordinate (sets of) activities or applications
 - Distributed
 - Running on different platforms using local coordinators
- Examples
 - Reach consistent agreement on the outcome of distributed transactions
 - Atomic transactions, 2PC
 - Coordinate auctioning activities
 - involves seller, auctioneer, buyers
 - Interactions between a customer and a supplier for ordering a product
 - request order, order goods, make payment



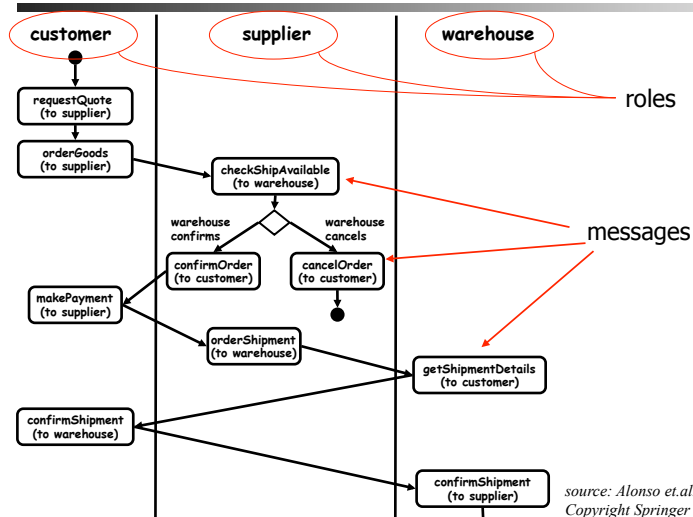
Conversations and Coordination Protocols

- Interactions form a **conversation**
 - sequences of operations (message exchanges)
 - maintain context information across invocations
- Interactions adhere to a **coordination protocol**
 - specifies a set of correct/accepted conversations
 - *vertical* protocols: specific to business area (e.g., product ordering protocol)
 - *horizontal* protocols: define common infrastructure (e.g., transactions)
- Different ways of modeling conversations
 - state machines
 - sequence diagrams
 - activity diagrams
- Middleware support can be provided, with various degrees of automation
 - conversation controllers
 - generic protocol handlers



© Prof.Dr.-Ing. Stefan DeBloch

Modeling Protocols - Activity Diagrams

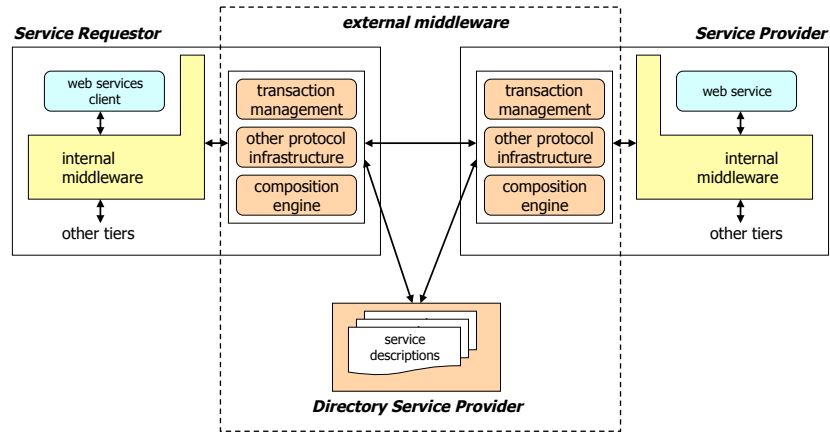


source: Alonso et.al.: Web Services, Springer, 2003
Copyright Springer Verlag Berlin Heidelberg 2003



© Prof.Dr.-Ing. Stefan DeBloch

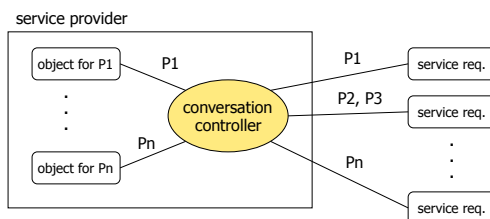
External Web Services Architecture



© Prof.Dr.-Ing. Stefan Deßloch

Conversation Controller

- Performs *conversation routing*
 - dispatch message to the appropriate "internal object"
 - one object for each instance of a conversation (e.g., an ordering session)
 - involves message correlation (conversation identifier), management of conversation context
 - example: session id
- Verifies *protocol compliance*
 - understand definition of the protocol (-> standardization of protocol descriptions)
 - check if all messages adhere to the protocol definition
- Can be implemented as a component of a SOAP router



© Prof.Dr.-Ing. Stefan Deßloch

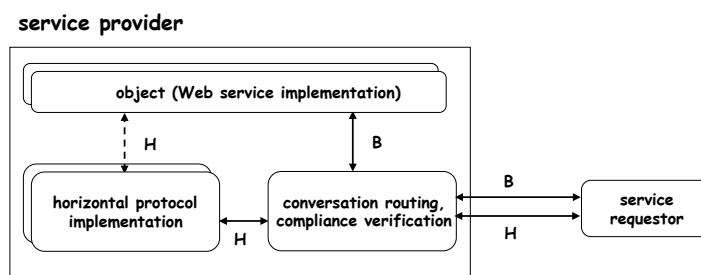
Generic Protocol Handlers

- Module that implements a specific coordination protocol
 - includes protocol-specific logic
 - processes and generates messages in accordance with the protocol rules
- Mostly applicable to horizontal protocols
 - example: transactions
- Forms of protocol execution support
 - handler realizes complete support, no intervention from the web service
 - Example: reliable messaging
 - handler and web service jointly realize the support
 - Example: atomic, distributed TAs
 - infrastructure coordinates sending/receiving prepare/commit/abort messages
 - web services decide over commit/abort, implement operations



© Prof.Dr.-Ing. Stefan DeBloch

Implementing Horizontal Protocols



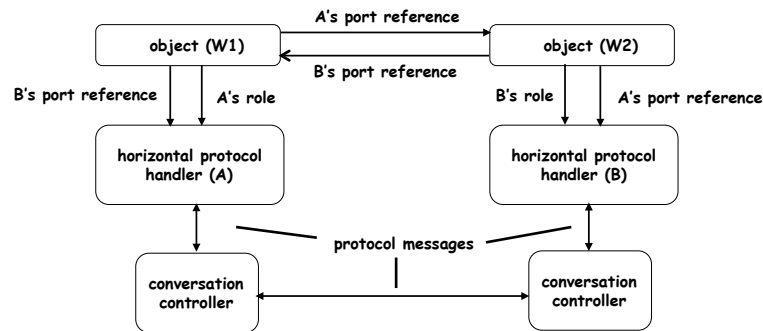
B: conversation compliant with a business protocol
H: conversation compliant with an horizontal protocol

source: Alonso et al.: *Web Services*, Springer, 2003
Copyright Springer Verlag Berlin Heidelberg 2003



© Prof.Dr.-Ing. Stefan DeBloch

Communicating Roles and Port References



source: Alonso et al.: Web Services, Springer, 2003
Copyright Springer Verlag Berlin Heidelberg 2003



© Prof.Dr.-Ing. Stefan DeBloch

Standardization

- Coordination infrastructure support for web services needs to be based on standards for
 - 1) generating and transporting unique conversation identifiers in SOAP headers
 - needed to map messages to conversations, and eventually to the objects handling them
 - 2) a framework and a set of (meta-) protocols for agreeing on which protocol is to be executed and how it is coordinated
 - 3) horizontal protocols
 - to separate horizontal protocol implementation from the individual web services
 - 4) protocol languages
 - to allow for protocol verification
- Web Services Coordination (WS-Coordination) Specification
 - standardizes 1), 2)
- Web Services Atomic Transaction (WS-AtomicTransaction) Specification
 - uses WS-Coordination framework to define coordination type for Atomic Transactions (i.e., it standardizes 3) for atomic TAs)
- Web Services Business Activity Framework (WS-BusinessActivity) Specification
 - same for (long-running) business transactions
- Standardized by the OASIS WS-TX technical committee
 - initial proposals by BEA, IBM, IONA, Microsoft



© Prof.Dr.-Ing. Stefan DeBloch

WS-Coordination

- Basic entities are **coordinators** and **participants** that wish to be coordinated
 - central coordination: all participants talk to a single coordinator
 - distributed coordination
 - each (or multiple) participant talks to its own coordinator
 - coordinators are chained together (subordinate coordinators act as participants)
- Abstractions to describe the interactions between coordinator and participants
 - **coordination protocol**
 - set of rules governing the conversation
 - example: 2PC
 - **coordination type**
 - set of logically related protocols
 - example: atomic transactions (completion, 2PC, volatile 2PC)
 - instance of a coordination type may involve several instances of the coordination protocols
- **Coordination context**
 - used to exchange coordination information among different parties
 - contains coordination type, identifier of the coordination type instance
 - placed within messages exchanged between parties (SOAP header)



© Prof.Dr.-Ing. Stefan DeBloch

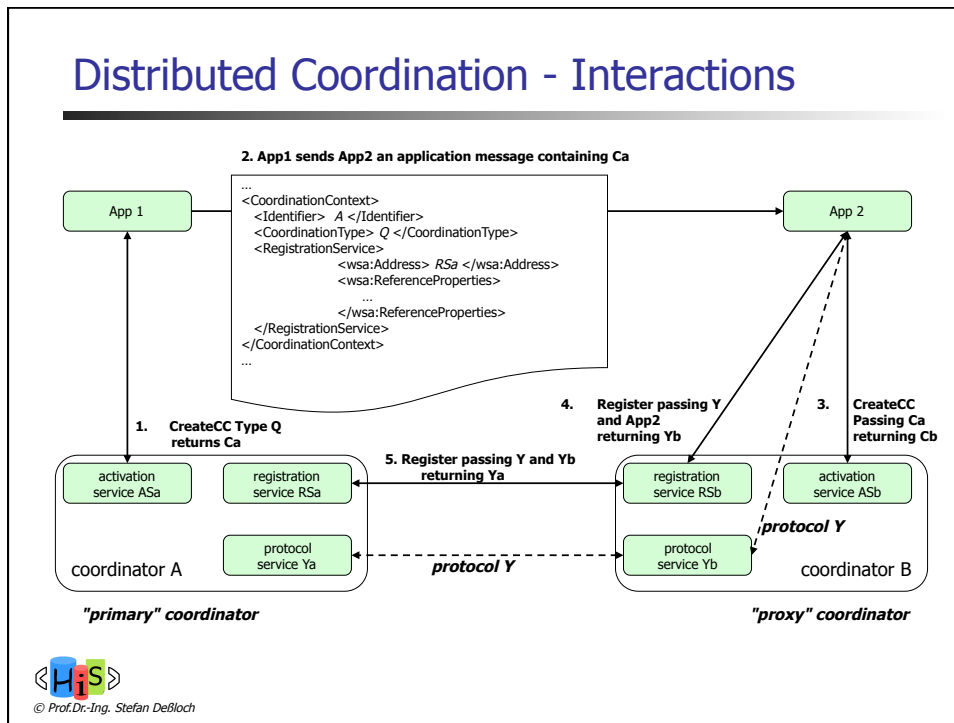
Coordinator/Participant Interactions

- Coordination service (coordinator) consists of
 - **Activation** service (generic)
 - Used by a participant to create coordination context (initiate instance of protocol type)
 - WS Interfaces: ActivationCoordinator, ActivationRequester
 - **Registration** service (generic)
 - Enable application to register for coordination protocols
 - provide endpoint information, role
 - WS Interfaces: RegistrationCoordinator, RegistrationRequester
 - (set of) **coordination protocols** (protocol-specific)
 - Specific to coordination type
- Extensibility
 - Publication of new coordination protocols
 - Definition of extension elements that can be added to protocols and messages



© Prof.Dr.-Ing. Stefan DeBloch

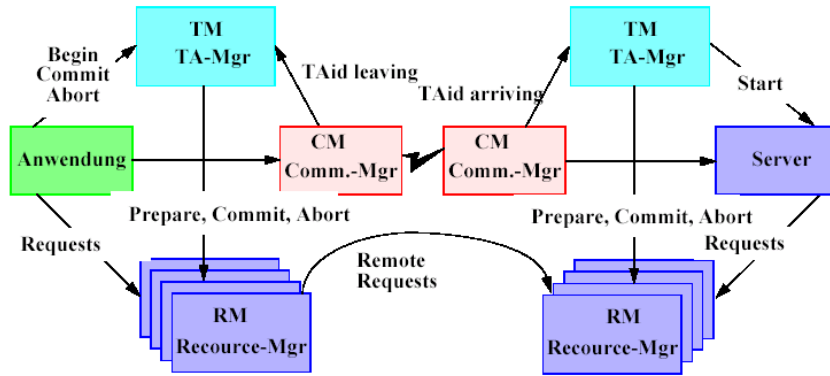
Distributed Coordination - Interactions



WS Atomic Transactions

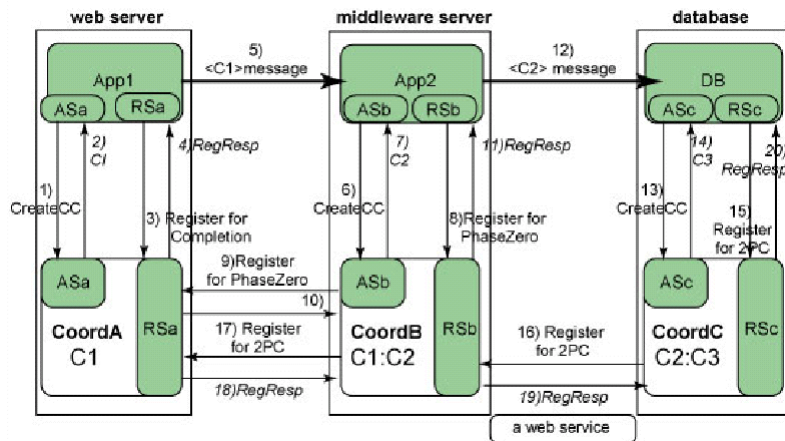
- Atomic Transactions (TA) coordination type
 - Defines type-specific commit protocols
 - Completion**: A participant (app creating the TA) registers so that it can tell the coordinator when/how to complete the TA (commit/abort)
 - 2PC**: a resource manager (RM) registers for this protocol to be included in the commit/abort decision
 - Hierarchical 2PC (local coordinators can be interposed as subordinate coordinators)
 - Two variants of 2PC
 - volatile 2PC**: a participant wants to be notified by the coordinator just before the 2PC begins
 - Example: participant caches, needs to communicate changes on cached data to DBMS before TA commits
 - durable 2PC**: a participant (e.g., DBMS) manages durable resources
 - Completion must be registered with the root coordinator
 - Participants can register for more than one protocol
 - Extension elements
 - Example: communicate **isolation levels**

X/Open DTP revisited ...



 © Prof.Dr.-Ing. Stefan DeBloch

AT WS-Coordination Flow



 © Prof.Dr.-Ing. Stefan DeBloch

AT WS-Coordination Flow (cont.)

- App1:
 - sends a **CreateCoordinationContext** message (1) to its local coordinator's Activation service ASa
 - create an atomic transaction T1
 - gets back in a **CreateCoordinationContextResponse** message (2) a **CoordinationContext** C1 containing the transaction identifier T1, the atomic transaction coordination type and CoordA's registration address RSa
 - sends a **Register** message (3) to RSa to register for the Completion protocol
 - gets back a **RegisterResponse** message (4), exchanging protocol service addresses for the coordinator and participant sides of the two-way protocol
 - sends an **application message** to App2 (5)
 - propagating the CoordinationContext C1 as a header in the message.
- App2:
 - decides to interpose local coordinator CoordB in front of CoordA
 - acts as a proxy to CoordA for App2
 - CoordA is the superior and CoordB is the subordinate
 - does this by sending a **CreateCoordinationContext** message (6) to the Activation service of CoordB (ASb) with C1 as input
 - getting back (7) a new **CoordinationContext** C2 that contains the same transaction identifier (T1) and coordination type, but has CoordB's registration address RSb.
 - **registers** with CoordB for the PhaseZero (volatile 2PC) protocol (8 and 11)
 - CoordB **registers** with CoordA for the PhaseZero protocol (9 and 10)
 - sends a **message to DB** (12), propagating CoordinationContext C2



© Prof.Dr.-Ing. Stefan Deßloch

AT WS-Coordination Flow (cont.)

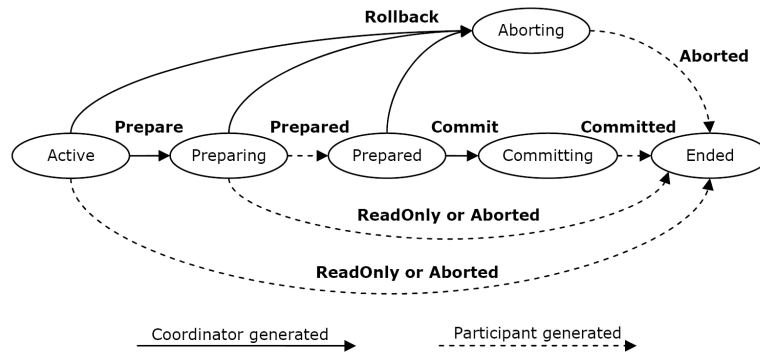
- DB:
 - decides to interpose its local coordinator CoordC by sending a **CreateCoordinationContext** message (13), further extending the superior-subordinate chain
 - gets back (14) a new **CoordinationContext** C3 that contains the same transaction identifier (T1) and coordination type, but CoordC's Registration service address RSc
 - **registers** with CoordC for the 2PC protocol because it is a resource manager (15 and 20)
 - causes CoordC to **register** with CoordB for the 2PC protocol (16 and 19)
 - causes CoordB to **register** with CoordA for the 2PC protocol (17 and 18)



© Prof.Dr.-Ing. Stefan Deßloch

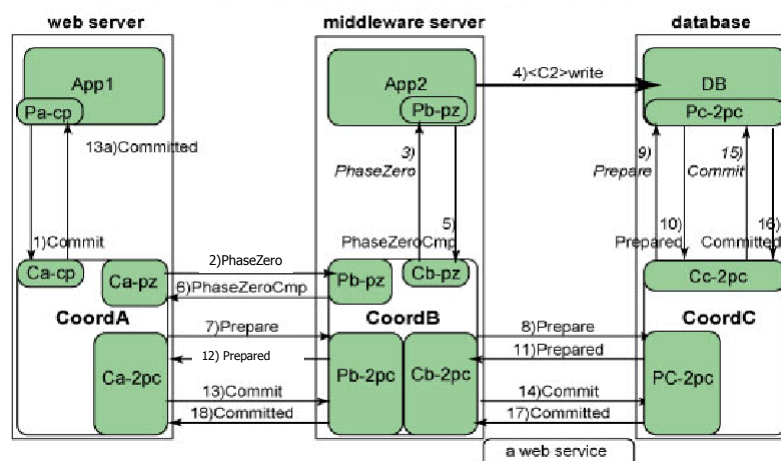
AT – 2PC Protocol

- Two-way protocol
 - Exchange of messages between coordinator and participant
- State Diagram
 - State reflects common knowledge of both parties




© Prof.Dr.-Ing. Stefan DeBloch

AT Coordination Protocol Flows




© Prof.Dr.-Ing. Stefan DeBloch

AT Coordination Protocol Flows (cont.)

- App1:
 - tries to **commit** the transaction using the Completion protocol (1)
- CoordA executes prepare-phase of **Volatile 2PC protocol**
 - has 1 participant registered for PhaseZero (CoordB), sends a **Prepare** message (2) to CoordB's PhaseZero Participant protocol service Pb-pz
 - CoordB relays **Prepare** message to App2 (3)
 - App2 sends its cached updates to DB
 - **application message** (4) propagates the CoordinationContext C2
 - sends a **Prepared** message (5) to CoordB
- CoordA executes prepare-phase of **durable 2PC protocol**
 - sends a **Prepare** message (7) to CoordB's 2PC Participant protocol service Pb-2pc
 - CoordB sends **Prepare** message (8) to CoordC's 2PC Participant protocol service Pc-2pc
 - CoordC tells DB to **Prepare** (9)
- CoordA **commits**
 - sends **Commit** message (13) to CoordB
 - Committed notification to App1 (13a) can also be sent
 - CoordB sends **Commit** message (14) to CoordC
 - CoordC tells DB to **commit T1**
 - DB receives the Commit message (15) and commits
 - **Committed** message returns (16, 17 and 18)



© Prof.Dr.-Ing. Stefan DeBloch

WS-BA – Business Activities Framework

- Characteristics (see discussion in chapter on WfMS)
 - Usually long-running
 - Responding to a request may take a long time
 - May consume lots of resources, perform a lot of work
 - Early commit of atomic subactivities/transactions
 - Forward recovery, compensation
- Goal: define protocols that "wrap" proprietary business activity mechanisms to achieve interoperability
- Design points
 - State transitions need to be reliably recorded
 - All request messages are acknowledged
 - Detect problems early
 - Response to a request is a separate operation
 - Not the output of the request
 - Avoid problems with timeouts of message I/O implementations



© Prof.Dr.-Ing. Stefan DeBloch

Compensation

- An action used to logically undo the effects of another action is called compensation action
 - Extends to real world actions
 - drilling a hole: throw away part
 - **Semantic Recovery**: Recovery schema based on compensation
 - Compensation very likely one of today's most frequently exploited techniques in transaction processing
- Compensation action is often dependent on context
 - E.g. writing an offer and sending it via mail to a customer
 - If letter is still in outbasket, simply remove it from outbasket
 - If letter is already received by the customer, write and send a countermanding letter
- Compensation often cannot recreate the same state that existed before the proper action had been performed
 - E.g. canceling a flight might cost a cancellation fee
 - Even more complicated, the cancellation fee might depend on the point in time, i.e. it is higher the later the cancellation is requested
- Compensation action may fail!



© Prof.Dr.-Ing. Stefan DeBloch

Sagas – Transactions and Compensation

- Sagas support specification of compensation actions in advance and run them automatically on abort
 - Sequence of (Sub-)Transaction/compensating action pairs
 - DBMS guarantees LIFO execution of compensation actions during abort/rollback of Saga
 - ACID for each sub-TA

Definition:

A **Saga** is a sequence $[(T_1, C_1), \dots, (T_n, C_n)]$ having the following properties:

1. T_1, \dots, T_n and C_1, \dots, C_n are two sets of transactions, such that C_i is the compensation function for T_i ,
2. $[(T_1, C_1), \dots, (T_n, C_n)]$ is executed as one of the following sequences:
 - i. $[T_1, \dots, T_n]$, if all T_i committed, or
 - ii. $[T_1, \dots, T_i, C_{i-1}, \dots, C_1]$ if T_i aborts and T_1, \dots, T_{i-1} committed before.



© Prof.Dr.-Ing. Stefan DeBloch

Business Activities Model

- Application is partitioned into **business activity scopes**
 - carries out business tasks using web services (participants)
 - mutually agreed outcome of all participants
- Participants registered with a coordinator of a BA
 - notify the coordinator about (successful) completion
 - may be asked by the coordinator to cancel an active task or to compensate a completed task
 - may indicate that it
 - cannot complete the task (and has cancelled it)
 - is leaving (exit) the BA (and has cancelled it)
 - has failed (during regular activities, when compensating or cancelling the task)
 - state of work is undetermined!
- Scopes may be arbitrarily nested



© Prof.Dr.-Ing. Stefan DeBloch

Business Activity (cont.)

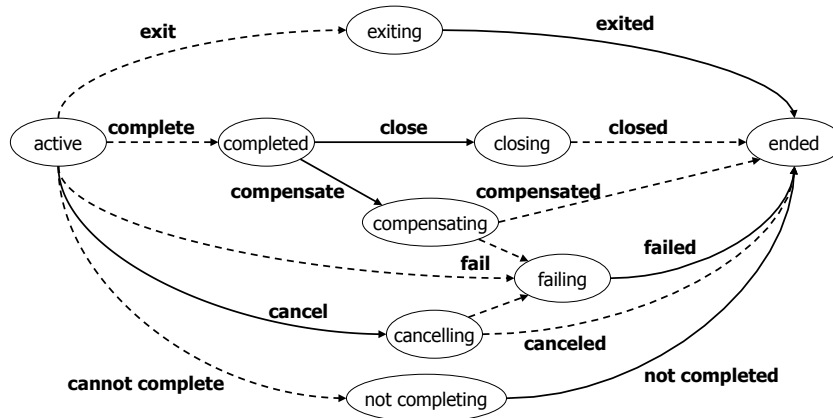
- Business Activity (BA) coordination types
 - AtomicOutcome: coordinator directs all participants to either close or compensate
 - MixedOutcome: coordinator may direct some participants to close, others to compensate
- BA protocol types
 - BusinessAgreementWithParticipantCompletion protocol
 - participant must know when it has completed all the work for a business activity
 - BusinessAgreementWithCoordinatorCompletion protocol
 - participant relies on coordinator to tell it when it has received all requests for work in the business activity



© Prof.Dr.-Ing. Stefan DeBloch

Business Agreement Protocol

BusinessAgreementWithParticipantCompletion – State Diagram



© Prof.Dr.-Ing. Stefan DeBiloch

Summary

- Coordination protocols
 - protocol defines set of correct conversations (WS message exchanges)
 - involves multiple partners, roles implemented as web services
 - vertical vs. horizontal protocols
 - different modeling approaches (e.g., activity diagrams)
- Infrastructure
 - conversation controller for internal routing based on conversation identifier
 - generic protocol handlers for horizontal protocols
- Coordination protocol infrastructure
 - WS-Coordination as a framework for supporting coordination protocols
 - central vs. distributed coordination
- WS-Transaction
 - based on WS-Coordination infrastructure
 - atomic transactions vs. business activities



© Prof.Dr.-Ing. Stefan DeBiloch