

IBM

IBM System z

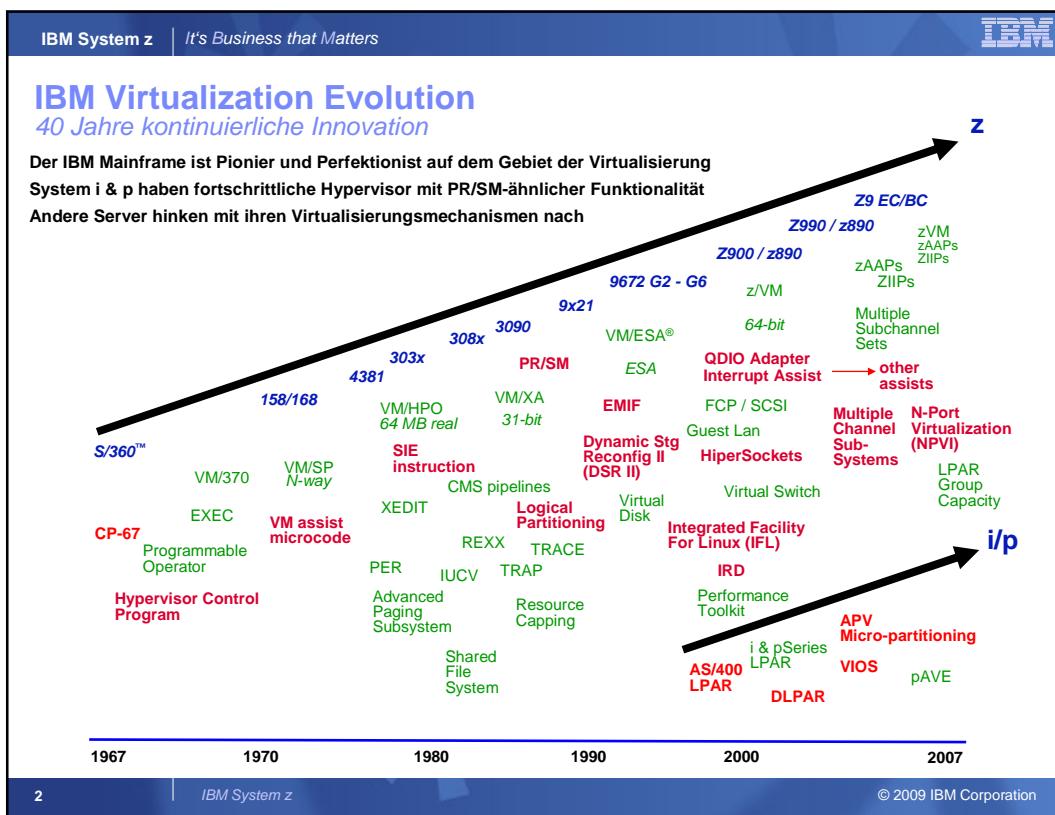
# The Mainframe and Linux

## Technischer Hintergrund von z/VM

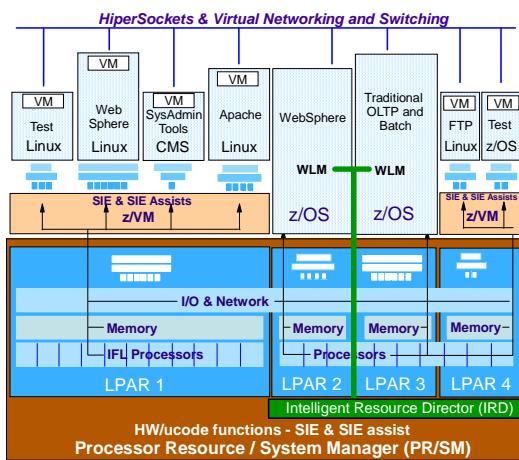
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## System z Virtualization Architecture

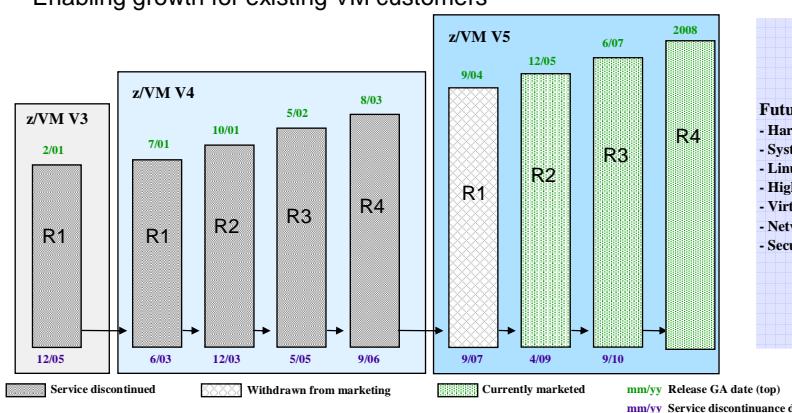


### Multi-Dimensionale Virtualisierung

- HW: PR/SM (LPAR's) & SW: zVM (VM's)
- Isolierte Umgebungen: EAL5 für PR/SM und EAL3+ für z/VM
- Hochentwickelter Hypervisor
- SIE Instruktion: Virtualisierung ist eingebaut, kein Add-On
- 10% der Integrierten Schaltkreise werden für Virtualisierung benutzt (SIE)
- Zeit- und Ereignisgesteuertes Dispatching
- SHARED ALL Architektur
- Jegliche virtuelle CPU kann auf jeder beliebigen physischen CPU betrieben werden. Sharing ist bis auf "1%" Ebene möglich
- Shared oder Dedizierter Pool an "CPUs"
- Garantierte LPAR Kapazität
- Physische & Virtuelle Ressourcen (CPU, I/O und Hauptspeicher) können dynamisch innerhalb und entlang der LPARs angepasst werden
- LPAR Zoning: Jede LPAR hat eine 0-Origin. Das erlaubt I/O-Zugriffe auf den Hauptspeicher einer LPAR ohne Hypervisor Eingriff
- z/VM kann virtuelle Devices erschaffen, die physisch nicht vorhanden sind
- z/VM hat raffinierte Scheduling-Algorithmen um das Gesamtsystem für Reaktion und Throughput zu optimieren

## System z: Evolution of z/VM

- z/VM Version 5: High-Value Virtualization Technology
  - Generating new business with Linux on System z
  - Enabling growth for existing VM customers



Future enhancements:

- Hardware and I/O
- Systems management
- Linux/VM synergies
- High availability
- Virtualization
- Networking
- Security

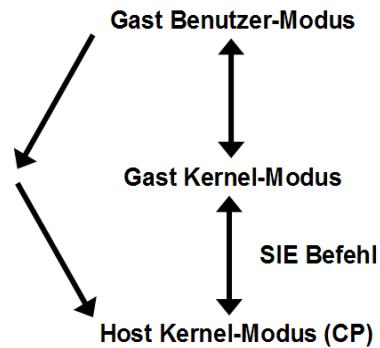
The EAL3+ evaluated copy of z/VM V5.1 is available as the Common Criteria Certification feature of z/VM V5.2

z/VM V5.2 will be withdrawn from marketing on June 15, 2007

Note: All statements regarding IBM's plans, directions, and intent are subject to change or withdrawal without notice, and represent goals and objectives only.

## System z Interpretive Execution –

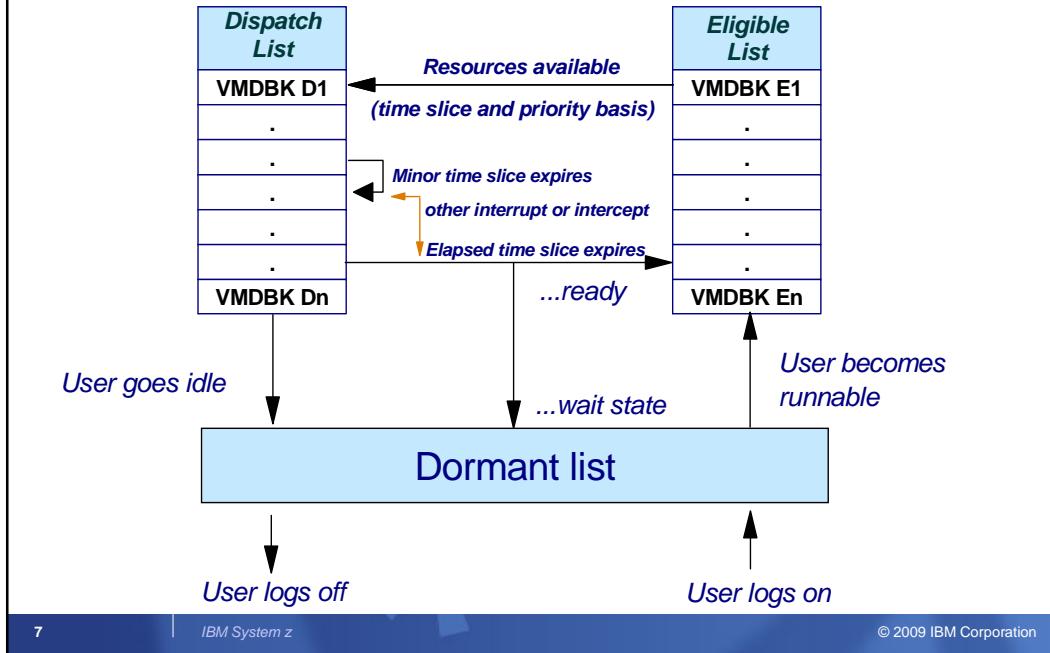
- Start Interpretive Execution (SIE) instruction
  - Establish the full architectural capabilities of an architecture for the guest
  - Supports MULTIPLE architectures CONCURRENTLY for multiple guests
  - Allows guests to run in NORMAL mode (no OS modifications needed/no TRAP'ing)
  - Reduces context switch time
  - Multiple control register sets
  - PR/SM AND z/VM exploit SIE  
No performance penalty for running z/VM in an LPAR



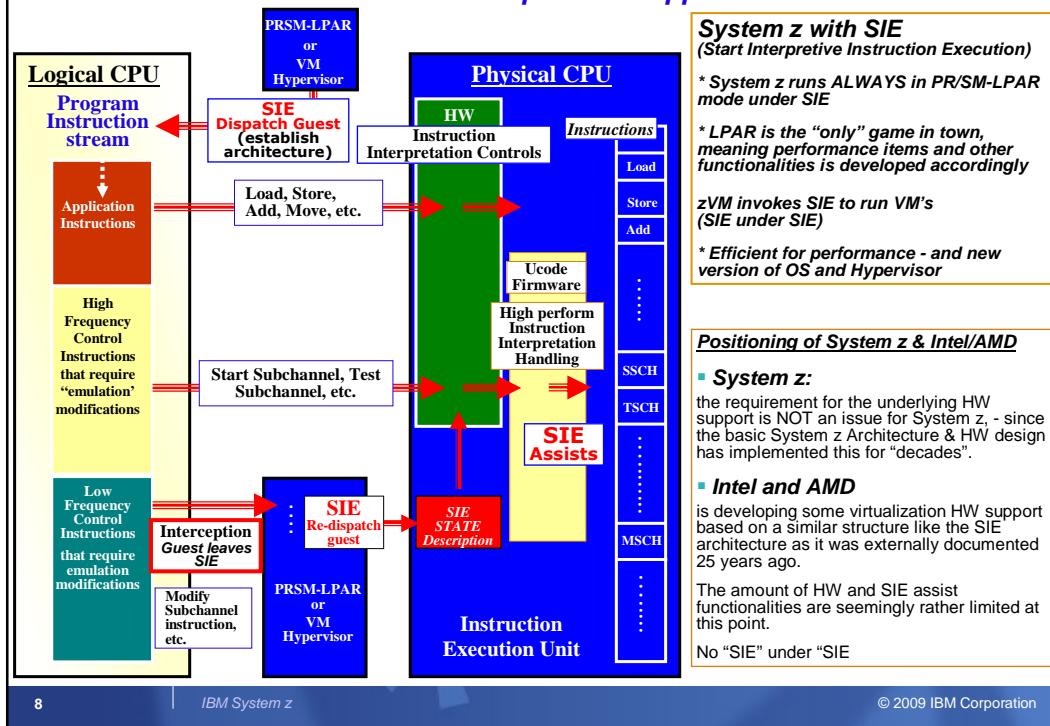
## How: Start Interpretive Execution (SIE)

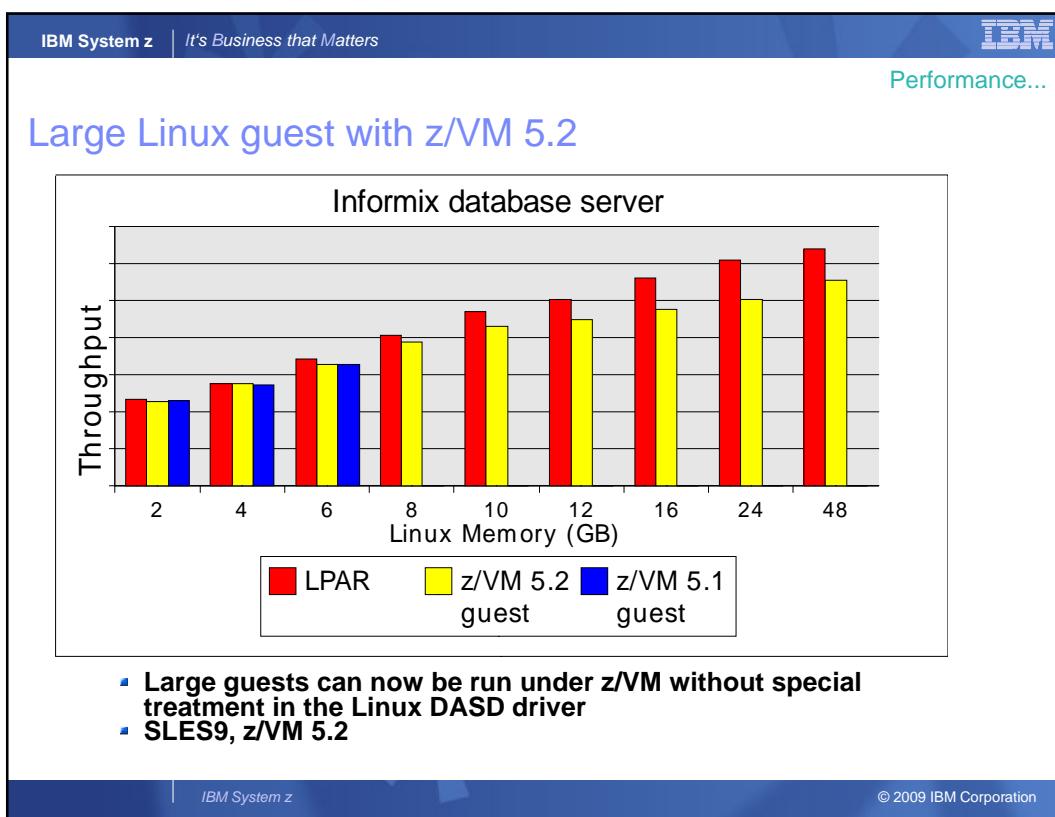
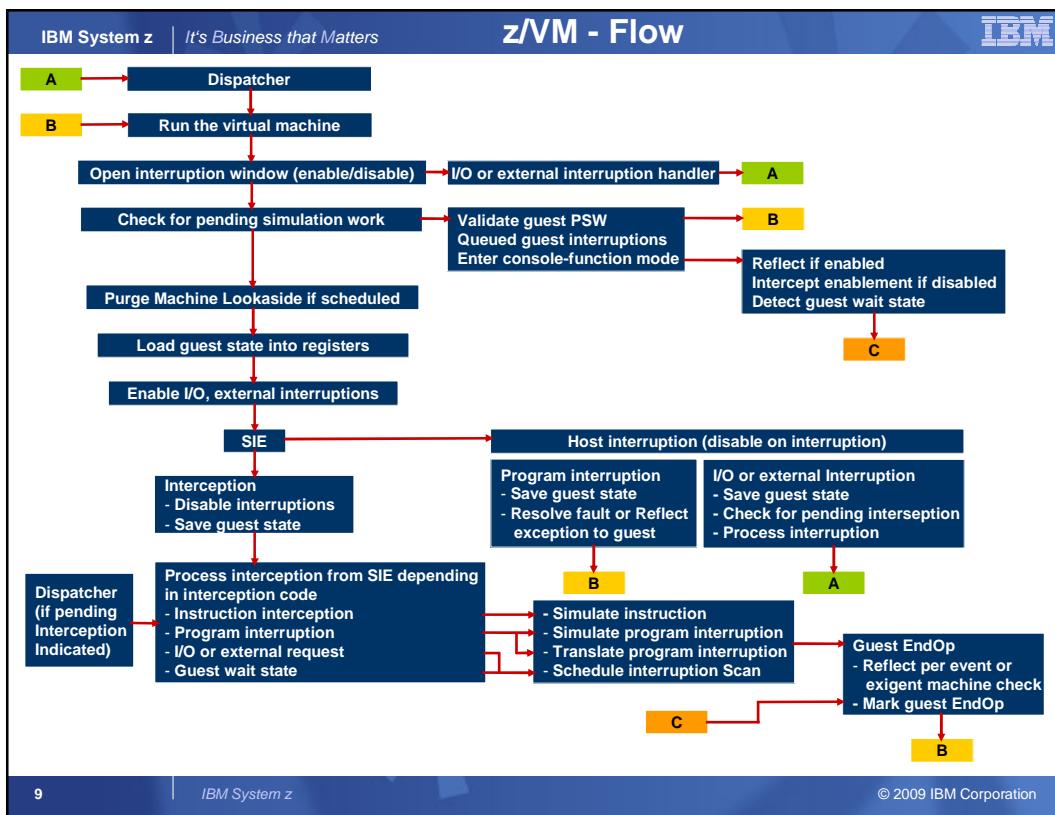
- SIE = “Start Interpretive Execution”, an instruction
- z/VM (like the LPAR hypervisor) uses the SIE instruction to “run” virtual processors for a given virtual machine.
- SIE has access to:
  - A control block that describes the virtual processor state (registers, etc.)
  - The Dynamic Address Translation (DAT) tables for the virtual machine
- z/VM gets control back from SIE for various reasons:
  - Page faults
  - I/O channel program translation
  - Privileged instructions (including CP system service calls)
  - CPU timer expiration (dispatch slice)
  - Other, including CP asking to get control for special cases
- CP can also shoulder tap SIE from another processor to remove virtual processor from SIE (perhaps to reflect an interrupt)

## Classic Scheduler / Dispatcher Picture

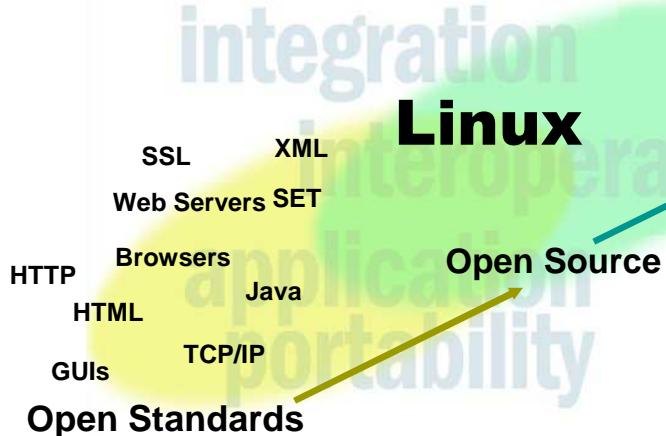


## Basic Direct HW virtualization – transparent to applications/OS





Questions?

# Thank You

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